

# Errors in Polynomial Interpolation

①

Suppose we want to approximate a (nice) function  $f(x)$  with a polynomial  $p(x)$  on ~~any~~ a fixed interval  $[a, b]$ .

We choose nodes  $x_0, \dots, x_n$  in  $[a, b]$ , and compute. Clearly  $f(x) = p(x)$  at  $x_0, \dots, x_n$ .

Let

$$E(n) = \max_{x \in [a, b]} |f(x) - p(x)|.$$

It is natural to expect  $E(n) \rightarrow 0$  as  $n \rightarrow \infty$ . In fact, for evenly spaced ~~nodes~~ nodes, it is often the case that  $E(n) \rightarrow \infty$  as  $n \rightarrow \infty$ !

Case study:  $f(x) = \frac{1}{1+x^2}$  on  $[-1,1]$ . ②

< chebyshev-nodes. nb >

We can give a few theorems about polynomial interpolation to shed some light on the matter.

Theorem. If  $p$  is the degree (at most)  $n$  polynomial interpolating  $f$  at  $n+1$  distinct nodes  $x_0, \dots, x_n$  in  $[a,b]$  and  $f^{(n+1)}$  is continuous, then for each  $x \in (a,b)$  there is some  $\xi \in (a,b)$  so

$$f(x) - p(x) = \frac{1}{(n+1)!} f^{(n+1)}(\xi) \prod_{i=0}^n (x-x_i).$$

③

Proof. If  $x$  is a node, we're done.

If not, let us fix  $x$  and set

$$\omega(t) = \prod_{i=0}^n (t - x_i)$$

$$c = \frac{f(x) - p(x)}{\omega(x)} \quad \leftarrow \text{this is a constant}$$

$$\varphi(t) = f(t) - p(t) - c\omega(t).$$

(Since  $x$  is not a node,  $\omega(x) \neq 0$ , so  $c$  exists.)

Now  $\varphi(t) = 0$  at all the nodes  $x_0, \dots, x_n$  and  $x$ .

We recall Rolle's theorem:

Theorem. Between any two roots of  $f$  there is a root of  $f'$ .

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The interval  $[a, b]$  now contains

$n+2$  roots of  $\varphi$

$n+1$  roots of  $\varphi'$

$n$  roots of  $\varphi''$

$\vdots$

$1$  root of  $\varphi^{(n+1)}$

At this point (call it  $\xi$ ), we have

$$0 = \varphi^{(n+1)}(\xi) = f^{(n+1)}(\xi) - p^{(n+1)}(\xi) + c \omega^{(n+1)}(\xi).$$

Now  $p$  is a polynomial of degree  $n$ , so  $p^{(n+1)}(x) \equiv 0$ . Now  $\omega(t)$  is a polynomial of degree  $n+1$  with leading term  $t^{n+1}$ , so  $\omega^{(n+1)}(t) \equiv (n+1)!$ . So

$$0 = f^{(n+1)}(\xi) + c(n+1)!$$

Substituting in the definition of  $c$ ...  $\square$

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Of course, this is not very helpful unless we can estimate  $\prod (x-x_i)$ .

Lemma. If  $x_i = a + ih$ ,  $i \in 0, \dots, n$ ,  $h = \frac{b-a}{n}$ , then for any  $x \in [a, b]$ ,

$$\prod_{i=0}^n |x-x_i| \leq \frac{1}{4} h^{n+1} n!$$

Proof. If  $x \in [x_j, x_{j+1}]$  we can show (homework) that

$$|x-x_j| |x-x_{j+1}| \leq \frac{h^2}{4}.$$

so

$$\prod_{i=0}^n |x-x_i| \leq \left( \prod_{i=0}^{j-1} (x-x_i) \right) \frac{h^2}{4} \left( \prod_{i=j+2}^n (x_i-x) \right)$$

§ Now for  $x_i \in [x_0, x_{j-1}]$ , we have  ~~$x_j > x_i$~~ ,  ~~$x_{j+1} > x_i$~~ ,

$x_{j+1} > x > x_j > x_i$ , so  $x-x_i < x_{j+1}-x_i$ .

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Similarly for  $x_i \in [x_{j+2}, x_n]$  we have

$$x_i > x_{j+1} > x > x_j, \text{ so } x_i - x < x_i - x_j.$$

So we have

$$\leq \left( \prod_{i=0}^{j-1} x_{j+1} - x_i \right) \frac{h^2}{4} \left( \prod_{i=j+2}^n x_i - x_j \right)$$

But  $x_{j+1} - x_i = (j+1-i)h$  and

$x_i - x_j = (i-j)h$ . So we really have

$$\leq h^j \cdot \frac{h^2}{4} \cdot h^{n-(j+2)+1} \cdot \prod_{i=0}^{j-1} (j-i)+1 \cdot \prod_{i=j+2}^n i-j$$

$$\leq \frac{h^{2+j+n-j-2+1}}{4} (j+1)! (n-j)!$$

$$\leq \frac{1}{4} h^{n+1} n!$$

where we use  $(j+1)! (n-j)! \leq n!$  for  $j \in 0, \dots, n$ . For  $j$  "in the middle" this is a gross overestimate but for  $j=0$  we

can do no better.  $\square$

⑦

We can now combine these to get

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Theorem. If  $f^{(n+1)}$  is continuous on  $[a, b]$  and bounded by  $M$  on  $[a, b]$ , then

$$|f(x) - p(x)| \leq \frac{1}{4(n+1)} M h^{n+1}.$$

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We can see from this that the key question is the growth of  $M$  as a function on  $n$ .

# Interpolation Errors and Divided Differences.

⑧

Last time we saw that in general for a polynomial we had

$$f(x) - p(x) = \frac{1}{(n+1)!} f^{(n+1)}(\xi) \prod_{i=0}^n (x-x_i)$$

Our mathematical examples show us that this error bound can blow up, along with the actual error.

<gridpoints\_and\_interpolation.nb>

Theorem. The Chebyshev nodes  $\hat{x}_i$  minimize

$$\max_{x \in [a, b]} \left| \prod (x - \hat{x}_i) \right| \text{ over all choices of } x_i.$$

$$x \in [a, b]$$

In fact, for Chebyshev nodes we have

$$|f(x) - P_n(x)| \leq \frac{1}{2^n (n+1)!} M, \text{ where } M \text{ bounds } f^{(n+1)}(x)$$



(note that this last bound only works on  $[-1, 1]$ ).

⑨

Now we show

Theorem. If  $p$  interpolates  $f$  on  $x_0, \dots, x_n$  then for any  $x$  which is not a node,

$$f(x) - p(x) = f[x_0, x_1, \dots, x_n, x] \prod_{i=0}^n (x - x_i).$$

~~Use~~  $\epsilon$

Proof. We know that if  $q$  interpolates

~~$f$~~  at  $x_0, \dots, x_n$  and  $\epsilon$  then

$$q(x) = p(x) + f_\bullet[x_0, \dots, x_n, \epsilon] \prod_{i=0}^n (x - x_i)$$

but at  $\epsilon = x$ ,  $q = f$ , so we have

$$f(x) - p(x) = f[x_0, \dots, x_n, x] \prod_{i=0}^n (x - x_i). \quad \square$$

Now this is weird and interesting because it suggests a connection between divided differences and derivatives:

Theorem. If  $f^{(n)}$  is continuous on  $[a, b]$ , for any (distinct)  $x_0, \dots, x_n$  in  $[a, b]$  there is some  $\xi$  in  $(a, b)$  so

$$f[x_0, \dots, x_n] = \frac{1}{n!} f^{(n)}(\xi).$$

Proof. Just combine these formulae for  $f(x) - p(x)$ .  $\square$

Corollary. If  $f$  is a polynomial of degree  $n$ , then  $f[x_0, \dots, x_i] = 0$  for  $i > n$ .

Example. Are the data points

x	1	-2	0	3	-1	7
y	-2	-56	-2	4	-16	376

formed by sampling a cubic?

(compute divided diffs.)